



Secure Xen on ARM: Status and Driver Domain Separation

Sang-bum Suh

sbuk.suh@samsung.com

SW Laboratories

Corporate Technology Operations

Samsung Electronics

Presented at Xen Summit Autumn 2007, Sun Microsystems

Contributors

- ◆ Sang-bum Suh
- ◆ Sung-min Lee
- ◆ Joo-young Hwang
- ◆ Jung-hyun Yu
- ◆ Sangdok Mo
- ◆ Chanju Park
- ◆ Bokdeuk Jeong
- ◆ Sungkwan Heo
- ◆ Jaemin Ryu
- ◆ Jun-young Sim
- ◆ Dong-hyuk Lee
- ◆ Igor Nabirushkin
- ◆ Alexander Trofimov
- ◆ Mikhail Levin
- ◆ Il-pyung Park
- ◆ Ho-soo Lee



Contents

◆ Overview and Status of Secure Xen on ARM Architecture 1.0

- ❖ Requirements for Beyond 3G Mobile Phone
- ❖ Goal and Architecture
- ❖ Development Environments
- ❖ Status of Secure Xen on ARM Architecture 1.0

◆ Driver Domain Separation: Architecture Exploration

- ❖ Motivation
- ❖ Driver Domain Separation: Architecture
- ❖ Summary
- ❖ Performance

◆ Future Work

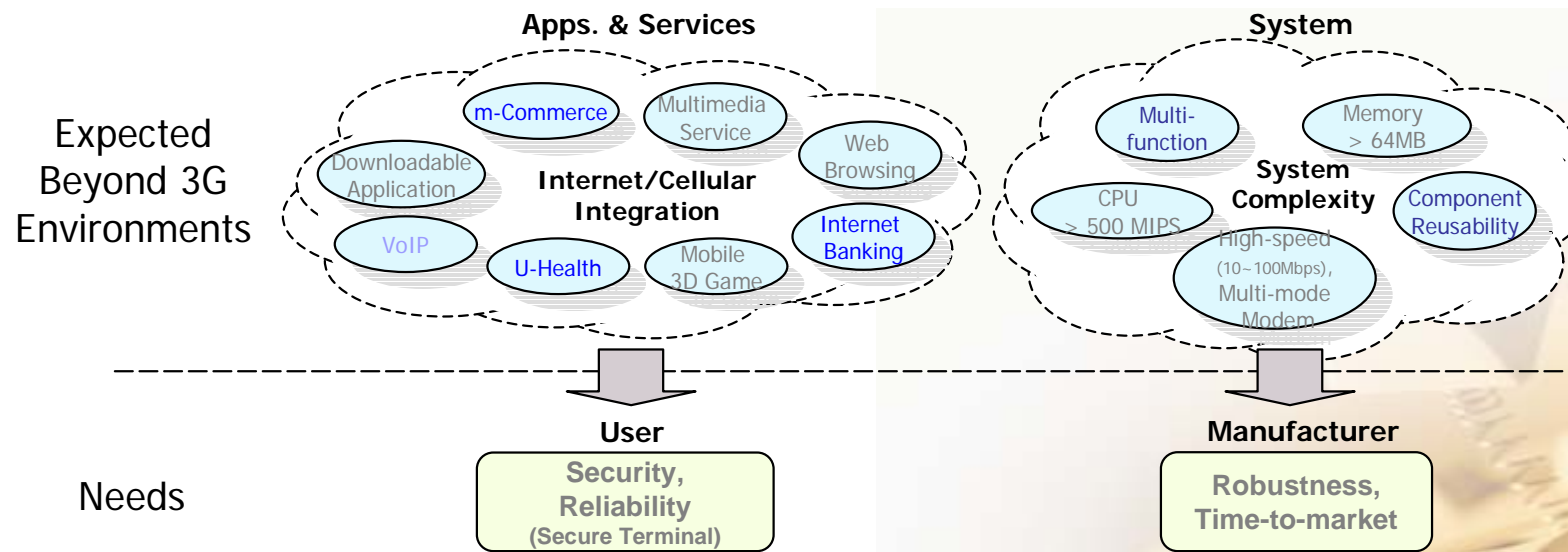


Overview and Status of Secure Xen on ARM Architecture 1.0



Requirements for Beyond 3G Mobile Phone

- ◆ End user: Secure and reliable mobile terminals for mobile Internet services using WiBro
- ◆ Manufacturer: Robustness though complexity of devices gets increased
- ◆ Contents provider: Protection of IP rights in end-user terminals
- ◆ Carrier companies: Open and Secure Mobile Platform
 - ❖ OSTI (Open Secure Terminal Initiative): NTT DoCoMo, Intel



Beyond 3G environments and Needs

Goal and Architecture

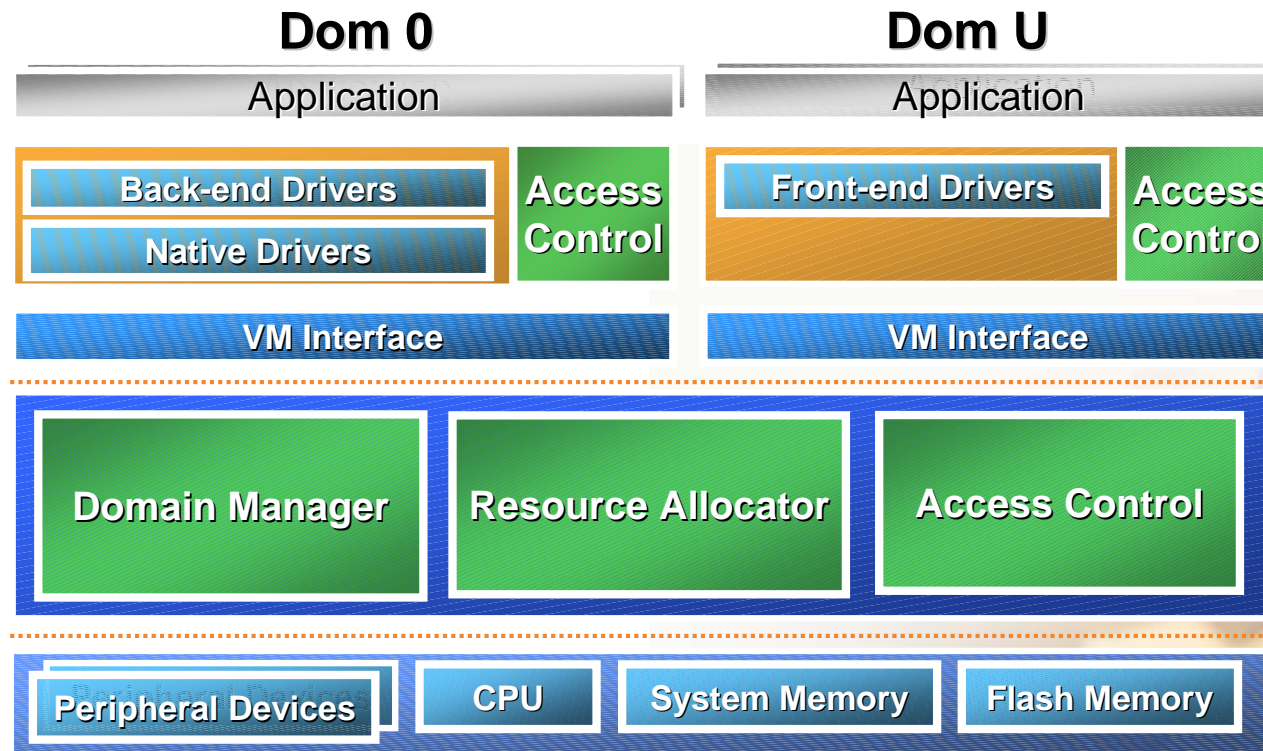
◆ Goal

- ❖ Light-weight secure virtualization technology for beyond 3G mobile phone

◆ Approach

- ❖ Design and implementation of

- VMM on ARM using Xen architecture: Xen on ARM
- Security features using Xen on ARM:
secure boot, secure SW installation, multi-layer fine-grained access control



Development Environments

◆ HW and SW Environments

❖ A Reference System for Implementation

➤ SW

- Xen : Xen-3.0.2
- Linux : ARM Linux-2.6.11
- GUI : Qtopia

➤ HW

- Processor : ARM-9 266Mhz (Freescale i.MX21)
- Memory : 64MB
- Flash : NOR 32MB / NAND 64MB
- LCD : 3.5 inch
- Network : CS8900A 10Base-T Ethernet Controller

❖ Development Environments

- OS : Fedora Core 6
- Cross-compiler: Montavista ARM GCC 3.3.1
- Debugger : Trace32 ICD (In Circuit Debugger)



Status of Secure Xen on ARM Architecture 1.0

◆ Xen on ARM:

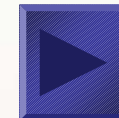
- ❖ Performance improved
- ❖ Video demo: game on Dom 0 and application/Qttopia on Dom U



video1

◆ Xen Security features:

- ❖ 5 access control modules and visualization supported:
 - Type Enforcement, Samsung proprietary, BiBA, Bell LaPadula, Chinese wall
 - GUI-based access control policy manager
- ❖ Video demo: access control mechanism against phishing attack



page21

◆ Driver domain separation: architecture exploration



Driver Domain Separation: Architecture Exploration



Motivation

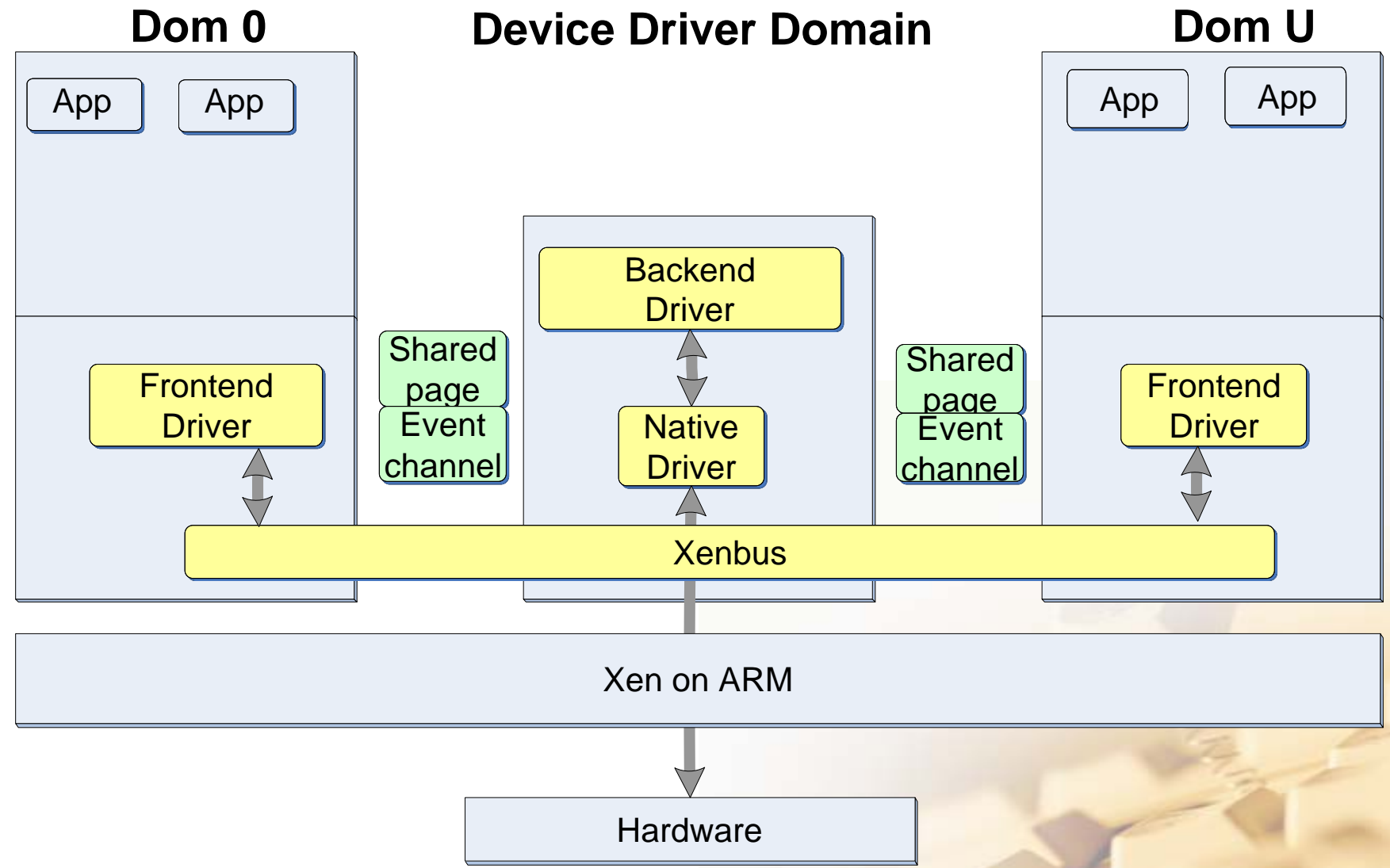
- ◆ Many downloadable services under beyond 3G mobile environments will be increased.
 - ❖ This requires an open mobile platform.

- ◆ Open platform will face problems with malware and bugs similar to PC.
 - ❖ Secure Xen can help an open mobile platform secure against malware.
 - ❖ However, bugs in device drivers may cause Dom 0 to stop working and the applications to have to restart.
 - Relatively short life cycle of peripheral chips in consumer electronics products.
 - Can test cases be updated quickly and be used to detect every bug during development ? Patch is likely.

- ☞ Device driver domain to be separated from Dom 0 (security applications running on Dom 0 in secure Xen on ARM) kernel.



Driver Domain Separation: Architecture



Summary

◆ Device driver domain

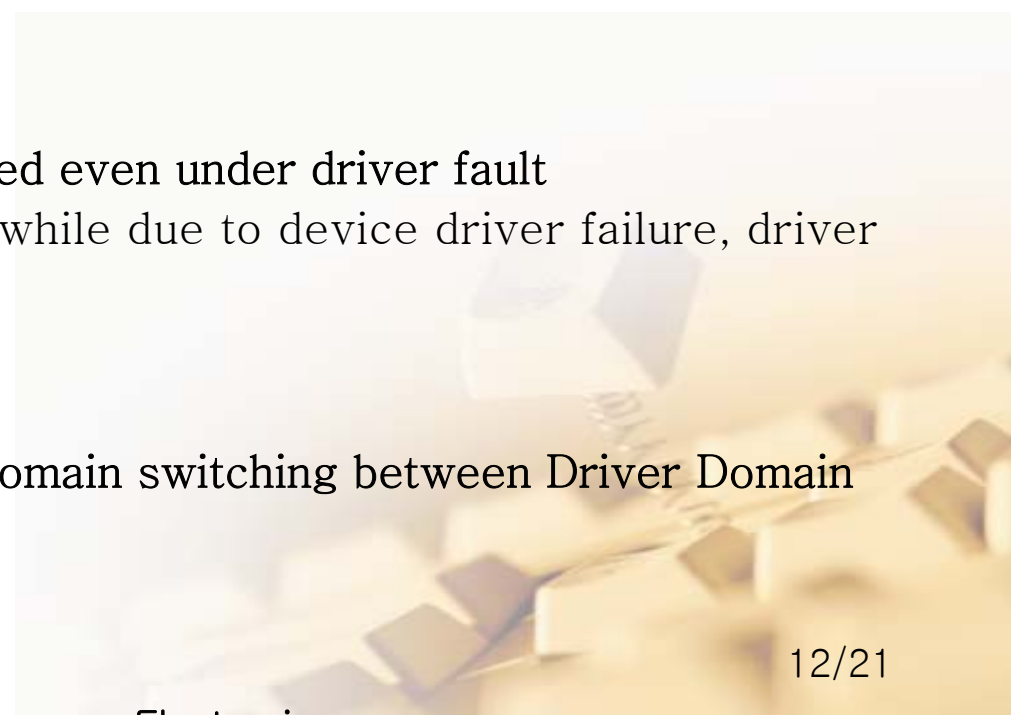
- ❖ Xen-Linux kernel, access control module, backend and native drivers
- ❖ Modification
 - RAMFS used for driver domain during booting
 - Xenbus, Xenstore, and Xen tools modified
 - Booting procedure modified
 - Booting Dom 0 => creating Device Driver Domain => initializing split device driver

◆ Advantage

- ❖ Service availability can be improved even under driver fault
 - Dom0 and Dom U can work, while due to device driver failure, driver domain has to be restarting.

◆ Disadvantage

- ❖ Performance degradation due to domain switching between Driver Domain and Dom 0

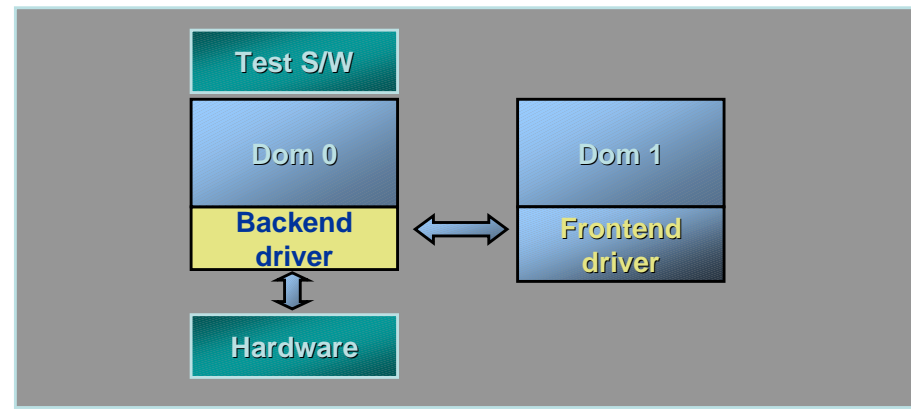


Performance (1/2)

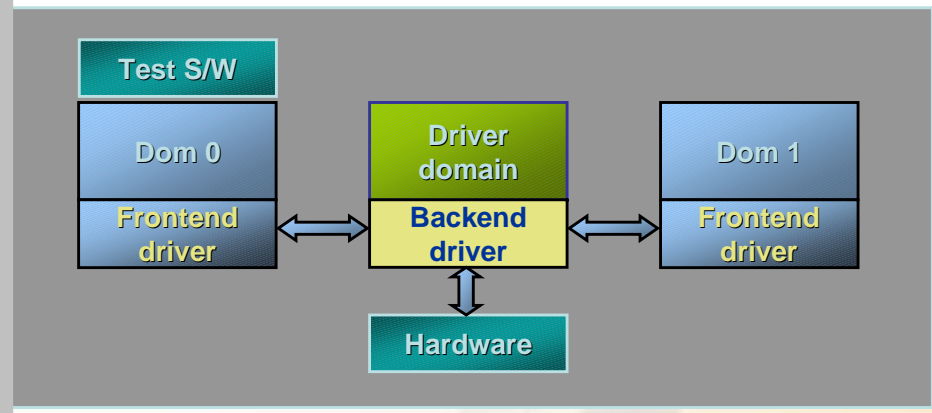
◆ Environments

- ❖ Virtual Network
- ❖ HW Platform: Freescale i.MX21
 - 266Mhz ARM926lrm5dmq
 - Memory: 64MB DDR
 - Network: CS8900A 10Base-T Ethernet Controller

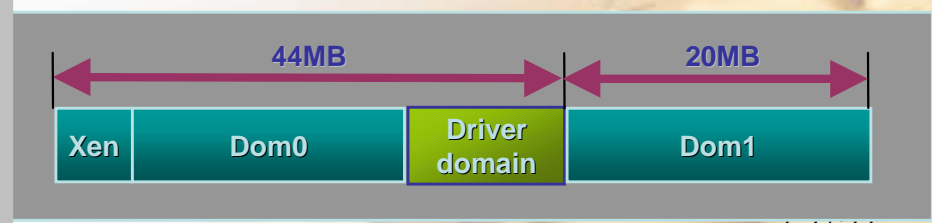
◆ Backend in dom0



◆ Backend in driver domain



◆ System memory configuration

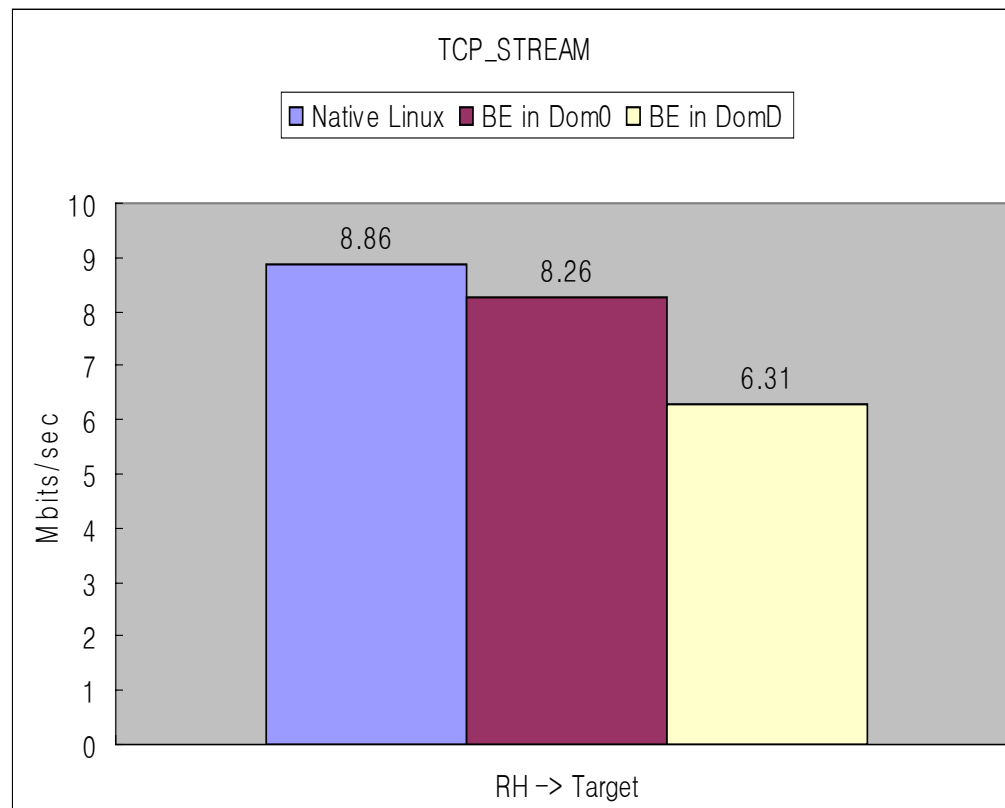


* Driver domain and Dom1 use Ramfs as a root file system. 13/21

Performance (2/2)

◆ Network Test: Netperf BMT

- ❖ Due to a problem with DMA of the HW, performance is degraded further.



- * **Native Linux:** network driver in native Linux
- * **BE in Dom0:** Backend driver in Dom0
- * **BE in DomD:** Backend driver in drive domain
- * **RH:** Remote Host

* **TCP_STREAM:** Measuring a bulk data transfer throughput

Future Work

- ◆ Performance improvement of driver domain separation
- ◆ Minimal OS kernel for driver domain
- ◆ State migration



Thank you for attention



Appendix

Access Control Module (1/2)

◆ Supporting 5 access control models

❖ Type Enforcement

- A classical access control model which can be enforced for comprehensive system resources protection
- Physical/virtual resources access control

❖ Proprietary

- Protecting a mobile device from resource drain attacks (e.g., CPU, memory, battery)

❖ Bell LaPadula

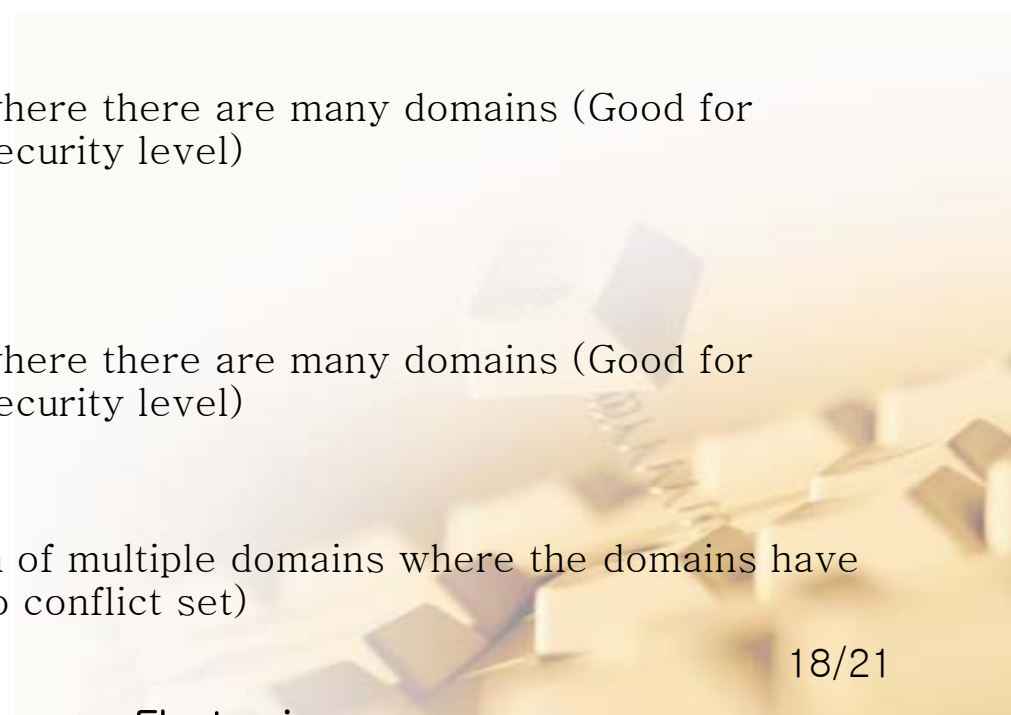
- Confidentiality model
- Virtual resources access control where there are many domains (Good for controlling information flow with security level)

❖ Biba

- Integrity model
- Virtual resources access control where there are many domains (Good for controlling information flow with security level)

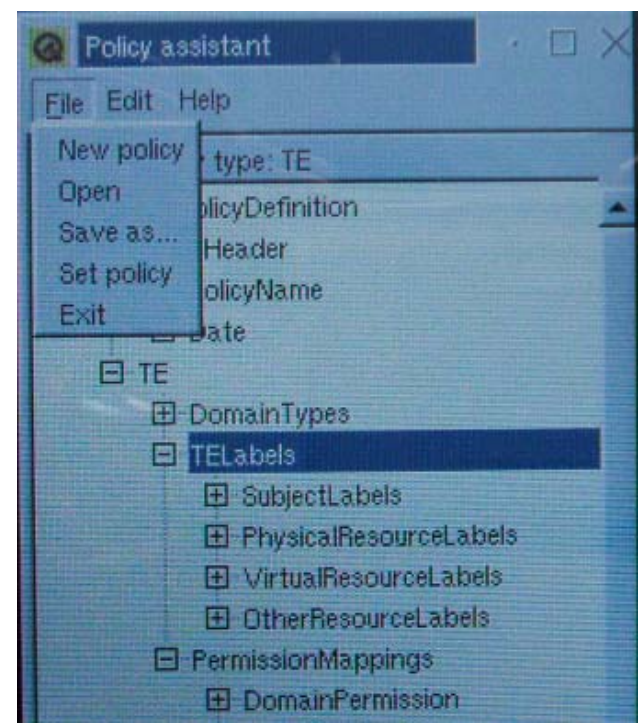
❖ Chinese Wall

- Preventing simultaneous execution of multiple domains where the domains have different interests (i.e., assigned to conflict set)



Access Control Module (2/2)

- ◆ GUI-based policy manager
 - ❖ Edits XML-based access control policies
 - ❖ Sets new access control policies dynamically



```
<?xml version="1.0" encoding="UTF-8" ?>
<!-- <SecurityPolicyDefinition xmlns="te_policy.xsd">
-->
- <SecurityPolicyDefinition>
+ <PolicyHeader>
- <TE>
  <!-- Definition of TE types which classify resources -->
+ <DomainTypes>
- <TELabels>
  <!-- Virtual machine labeling -->
- <SubjectLabels>
+ <VMLabel>
+ <VMLabel>
+ <VMLabel>
</SubjectLabels>
  <!-- Physical resource labeling -->
+ <PhysicalResourceLabels>
- <PhyResourceLabel>
  <Type>te_KPP</Type>
  <PIRQLabel>KPP</PIRQLabel>
  <IOMEMLabel>KPP</IOMEMLabel>
</PhyResourceLabel>
- <PhyResourceLabel>
  <Type>te_Flash</Type>
  <IOMEMLabel>FLASHMEM</IOMEMLabel>
</PhyResourceLabel>
- <PhyResourceLabel>
  <Type>te_CONSOLE</Type>
  <PIRQLabel>UART1</PIRQLabel>
  <IOMEMLabel>UART1</IOMEMLabel>
</PhyResourceLabel>
</PhysicalResourceLabels>
```

Example of the XML-based TE policy

Secure SW Installation

- ◆ Basic assumptions about software on the secure domain
 - ❖ A small set of software (not much) can be installed by only trusted parties (i.e., manufacturer or service providers verified by the manufacturer)
 - ❖ The trusted parties must rigorously test the software based on advanced quality assurance methodology during the development phase
- ◆ Secure SW installer installs only software digitally signed by a manufacturer
- ◆ Access control at the secure domain (Dom0) allows only authentic secure SW installer to create executable files on the domain
 - ❖ Even in case a device owner downloads or creates files on the secure domain, they cannot be executed



Demonstration Scenario



◆ Connecting to a phishing site

- ❖ Alice connects to a phishing server with her mobile phone after receiving an email fraudulently saying launch of UCC services from her favorite web site
- ❖ She downloads and installs malware masqueraded as genuine SW from that site

◆ With a conventional single OS-based mobile phone

- ❖ Malware corrupts kernel and sends her sensitive information to an attacker while she is using the Internet banking service



video2

◆ With a secure Xen-based mobile phone (with secure domain and normal domain)

- ❖ Even in case malware corrupts kernel of the normal domain, there is no information leakage or availability threat owing to domain separation and mandatory access control
- ❖ Secure SW installer installs Gifviewer signed by a manufacturer successfully but fails to install Pacman whose digital signature is invalid
 - Assumption: communication channel between the secure SW installer and manufacturer site which provides downloadable SW is encrypted



video3

